

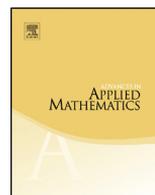


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New characterisations of tree-based networks and proximity measures



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ABSTRACT

Phylogenetic networks are a type of directed acyclic graph that represent how a set X of present-day species are descended from a common ancestor by processes of speciation and reticulate evolution. In the absence of reticulate evolution, such networks are simply phylogenetic (evolutionary) trees. Moreover, phylogenetic networks that are not trees can sometimes be represented as phylogenetic trees with additional directed edges placed between their edges. Such networks are called *tree-based*, and the class of phylogenetic networks that are tree-based has recently been characterised. In this paper, we establish a number of new characterisations of tree-based networks in terms of path partitions and antichains (in the spirit of Dilworth's theorem), as well as via matchings in a bipartite graph. We also show that a temporal network is tree-based if and only if it satisfies an antichain-to-leaf condition. In the second part of the paper, we define three indices that measure the extent to which an arbitrary phylogenetic network deviates from being tree-based. We describe how

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these three indices can be computed efficiently using classical results concerning maximum-sized matchings in bipartite graphs.

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1. Introduction

Phylogenetic networks are of increasing interest in the literature as they allow for the representation of reticulate (non-tree-like) processes in evolution. From a mathematical perspective, there are several particularly attractive classes of phylogenetic networks. One of those classes is tree-based networks. Intuitively, a phylogenetic network is tree-based if it can be obtained from a phylogenetic tree \mathcal{T} by simply adding edges whose end-vertices subdivide edges of \mathcal{T} . Formalised and studied in [5], tree-based networks have since been studied in a number of recent papers [1,7,9,11,12], in a variety of contexts.

In this paper, we establish several new characterisations of tree-based networks. These characterisations are based on antichains, path partitions, and matchings in bipartite graphs, and complement the previous characterisations based on bipartite matchings [9, 12]. Furthermore, with the aid of these characterisations, we explore indices quantifying the closeness of an arbitrary phylogenetic network \mathcal{N} to being tree-based. Each of the considered indices is computable in time polynomial in the size of \mathcal{N} , that is, in the number of vertices in \mathcal{N} , by finding a maximum-size matching in certain bipartite graphs based on \mathcal{N} .

The paper is organised as follows. The rest of the introduction contains some formal definitions and previous results. In Section 2, we state the new characterisations of tree-based networks as well as a characterisation of tree-based networks within the class of temporal networks. Except for one characterisation which is proved in Section 4, the proofs of these characterisations are given in Section 3. In Section 4, we consider three indices quantifying the extent to which an arbitrary phylogenetic network is tree-based. We end the paper with a brief conclusion and further questions in Section 5.

1.1. Definitions

Throughout the paper, X denotes a non-empty finite set. A *phylogenetic network* \mathcal{N} on X is a rooted acyclic digraph with no edges in parallel and satisfying the following properties:

- (i) the (unique) root has out-degree two;
- (ii) a vertex with out-degree zero has in-degree one, and the set of vertices with out-degree zero is X ; and
- (iii) all other vertices either have in-degree one and out-degree two, or in-degree two and out-degree one.

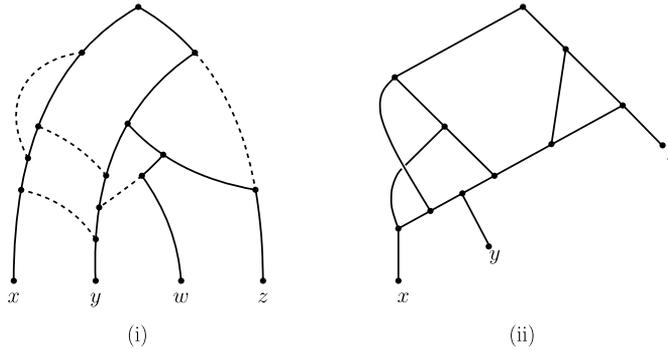


Fig. 1. Examples of (i) a tree-based network, showing a base tree (solid edges), and (ii) a non-tree-based network.

For technical reasons, if $|X| = 1$, we additionally allow \mathcal{N} to consist of the single vertex in X . The vertices of \mathcal{N} with out-degree zero are referred to as *leaves*. Furthermore, the vertices of \mathcal{N} with in-degree two and out-degree one are *reticulations*, while the root together with the vertices with in-degree one and out-degree two are *tree vertices*. The edges directed into a reticulation are *reticulation edges*; all other edges are *tree edges*. In the literature, a phylogenetic network as defined here is sometimes called a *binary phylogenetic network*. A *rooted (binary) phylogenetic X-tree* is a phylogenetic network on X with no reticulations.

A phylogenetic network $\mathcal{N} = (V, E)$ on X is a *tree-based network* if \mathcal{N} has a rooted spanning tree (V, E') , where $E' \subseteq E$, that has all its leaves in X . We refer to this spanning tree as a *base tree* for \mathcal{N} . The initial study of tree-based networks [5] included the presentation of a polynomial-time algorithm (based on a characterisation involving 2-SAT) for determining whether or not an arbitrary phylogenetic network is tree-based and, if so, constructing a base tree for it. Fig. 1 shows an example of a phylogenetic network that is tree-based, and one that is not. As with all drawings of phylogenetic networks in this paper, edges are directed down the page and, thus, away from the root.

Lastly, in various places in the paper, we use the following operation. Let G be a directed graph and let e be an edge of G . The operation of deleting e , adding a new vertex x , and adjoining x to each of the end-vertices of e (orienting the new edges in the original direction of e) is to *subdivide* e . Any directed graph obtained from G by a sequence of edge subdivisions is called a *subdivision* of G .

1.2. Previous results

Let \mathcal{N} be a phylogenetic network on X . Let T be the set of tree vertices in \mathcal{N} that are parents of a reticulation and let R be the set of reticulations in \mathcal{N} . Zhang [12] defined the following bipartite graph to characterise tree-based networks. Let $\mathcal{Z}_{\mathcal{N}}$ be the graph with vertex set $T \cup R$ and edge set

$$\{\{t, r\} : t \in T, r \in R, \text{ and } (t, r) \text{ is an edge in } \mathcal{N}\}.$$

In particular, Zhang [12] established the following characterisations.

Theorem 1.1. *Let \mathcal{N} be a phylogenetic network. Then the following are equivalent:*

- (i) \mathcal{N} is tree-based.
- (ii) The bipartite graph $\mathcal{Z}_{\mathcal{N}}$ has a matching such that each reticulation is matched.
- (iii) The bipartite graph $\mathcal{Z}_{\mathcal{N}}$ has no maximal path that starts and ends with reticulations.

The proofs of our main results make use of these characterisations. It is clear that (ii) implies (iii), while the reverse implication relies on Hall's theorem [6] (see [3, Theorem 16.4]) on matchings in bipartite graphs.

There are several other characterisations of what it means for a phylogenetic network $\mathcal{N} = (V, E)$ to be a tree-based network:

- (i) The existence of an 'admissible' subset of the edges of \mathcal{N} [5].
- (ii) The existence of an independent subset of edges E' of E for which $(V, E - E')$ is a base tree of \mathcal{N} [5].
- (iii) The existence of a matching in a certain bipartite graph whose vertices consist of 'omnians' (vertices whose only children are reticulations) and reticulations [9].

While we mention these characterisations, they will play no further part in the paper.

2. New characterisations

An *antichain* in a directed graph is a subset S of vertices with the property that, for all distinct $u, v \in S$, there is no directed path from u to v . Any tree-based network \mathcal{N} satisfies the *antichain-to-leaf property*, which says that for any antichain of k vertices, there exist k vertex disjoint paths from the elements of the antichain to the leaves of \mathcal{N} [5]. However, this property is not sufficient to ensure a network is tree-based, as the counterexample in Fig. 1(ii) illustrates. Note that if an arbitrary phylogenetic network satisfies the antichain-to-leaf property, then it satisfies the corresponding property for *edge* disjoint paths, and in fact is equivalent to it via an application of Menger's theorem [10] (see [3, Theorem 7.16]).

The main result of this section is the following theorem. This theorem provides five properties of a phylogenetic network $\mathcal{N} = (V, E)$ that are equivalent to being tree-based. Four of these properties, (II)–(V), can be viewed as providing different ways to strengthen the antichain-to-leaf property. The final property, (VI), provides another characterisation in terms of bipartite graphs. Let $\mathcal{G}_{\mathcal{N}}$ denote the bipartite graph whose vertex bipartition is $\{V_1, V_2\}$, where each of V_1 and V_2 is a copy of V , and with an edge joining a vertex

$u \in V_1$ and a vertex $v \in V_2$ precisely if (u, v) is an edge in \mathcal{N} . To illustrate, consider the phylogenetic network \mathcal{N} shown in Fig. 3(i). The bipartite graph $\mathcal{G}_{\mathcal{N}}$ is shown in Fig. 3(ii).

Theorem 2.1. *Let $\mathcal{N} = (V, E)$ be a phylogenetic network on X . The following are equivalent:*

- (I) \mathcal{N} is tree-based;
- (II) \mathcal{N} has an antichain $\mathcal{A} \subseteq V$, and a partition Π of V into $|\mathcal{A}|$ chains each of which forms a path in \mathcal{N} ending at a leaf in X ;
- (III) For all $U \subseteq V$, there exists a set of vertex disjoint paths in \mathcal{N} each ending at a leaf in X such that each element of U is on exactly one path;
- (IV) There is no pair of subsets $U_1, U_2 \subseteq V$ such that $|U_1| > |U_2|$ and
 - (i) every path from a vertex in U_1 to a vertex in X traverses a vertex in U_2 , and
 - (ii) for $\{i, j\} = \{1, 2\}$, if there is a path from a vertex in U_i to a vertex in U_i , then this path traverses a vertex in U_j .
- (V) The vertex set of \mathcal{N} can be partitioned into a set of vertex disjoint paths, each of which ends at a leaf in X .
- (VI) The bipartite graph $\mathcal{G}_{\mathcal{N}}$ has a matching of size $|V| - |X|$.

Remarks:

- Property (II) is related to a classical result in combinatorics, namely Dilworth’s theorem [4]. A version of this theorem states that every finite poset P has an antichain A and a partition Π of P into $|A|$ chains (see [3, Theorem 19.5]). If we regard a network as a poset in the usual way, the additional requirement in (II) is that the chains must form paths; one cannot simply ‘jump over’ other vertices. It is instructive to see why this fails for the counterexample to the ‘antichain-to-leaf property’ in [5], reproduced in Fig. 1(ii). This network has a maximum-sized antichain of size three (e.g. $\{x, y, z\}$), and so the vertex set can be partitioned into three chains (by Dilworth’s theorem). However, the network cannot be partitioned into three paths. A partition into four paths is shown in Fig. 2.
- Notice that the condition in (II) forces $|\mathcal{A}| = |X|$, since each element x in X has to be in exactly one path from Π . But, \mathcal{A} need not necessarily be equal to X , indeed one could have $\mathcal{A} \cap X = \emptyset$.
- Property (VI) complements the two existing characterisations of tree-based networks via matchings [9,12].

A second result in this section is a characterisation of tree-based networks within the class of ‘temporal’ networks. Introduced in [2], a phylogenetic network $\mathcal{N} = (V, E)$ is *temporal* if there is a map $\lambda : V \rightarrow \mathbb{R}$ so that $\lambda(u) < \lambda(v)$ for each tree edge (u, v) , and $\lambda(u) = \lambda(v)$ for each reticulation edge (u, v) , in which case λ is a *temporal map* for \mathcal{N} . As illustrated above, the antichain-to-leaf property is a necessary but not sufficient

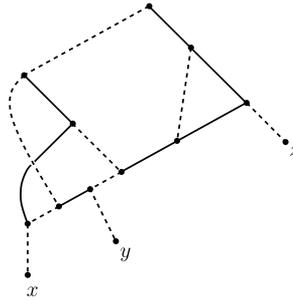


Fig. 2. For the non-tree-based network from Fig. 1(ii), the vertex set can be partitioned into four paths (dashed), but not into three (the size of the maximum-sized antichain).

condition for a phylogenetic network to be tree-based. However, within the class of temporal networks, it is sufficient.

Theorem 2.2. *Let \mathcal{N} be a temporal network. Then \mathcal{N} is tree-based if and only if \mathcal{N} satisfies the antichain-to-leaf property.*

3. Proof of Theorems 2.1 and 2.2

In this section, we establish the equivalence of (I)–(V) in Theorem 2.1 as well as Theorem 2.2. The proof of the equivalence of (I) and (VI) in Theorem 2.1 is done independently as Corollary 4.2. We begin with a lemma.

Lemma 3.1. *Let T be a subdivision of a rooted binary tree with vertex set V . Then the following property holds:*

(P) *For any non-empty subset U of V there exists a set of vertex disjoint (directed) paths in T each of which ends at a leaf of T and each vertex in U lies on exactly one path.*

Proof. We apply induction on the number n of vertices of T . For $n = 1$, (P) trivially holds. Suppose that $n \geq 2$ and (P) holds for all subdivisions of a rooted binary tree with at most $n - 1$ vertices. Let U be an arbitrary subset of vertices of T . Since $n \geq 2$, it follows that T either has

- (i) a leaf x whose parent, u say, has degree 2, or
- (ii) a vertex v that is a parent of two leaves, x and y say.

We establish the induction hypothesis in each case, starting with (i).

For (i), let T' be the subdivision of a rooted binary tree obtained from T by deleting x and its incident edge, so that u is now a leaf of T' . Let

$$U' = \begin{cases} U, & \text{if } U \text{ does not contain } x; \\ U - \{x\}, & \text{if } U \text{ contains } x \text{ and also contains } u; \\ (U - \{x\}) \cup \{u\}, & \text{if } U \text{ contains } x \text{ but not } u. \end{cases}$$

Observe that U' is a subset of vertices of T' . Therefore, as T' has $n - 1$ vertices, it follows by induction that (P) holds and so there is a set of disjoint paths in T' each of which ends at a leaf of T' and each vertex in U lies on exactly one path. Now one of these paths ends at u . Replacing this path with the one that extends it to end at x gives a set of vertex disjoint paths in T , each of which ends at a leaf of T , and each vertex in U lies on exactly one path. Thus the lemma holds for (i).

Now consider (ii). Let T' be the subdivision of a rooted binary tree obtained from T by deleting y and its incident edge. Note that T' has $n - 1$ vertices. If U does not contain y , then let $U' = U$. By induction, there is a set of vertex disjoint paths in T' each of which ends at a leaf of T' and each vertex in U' lies on exactly one path. This set of paths also works for U in T . On the other hand, if U does contain y , then let $U' = U - \{y\}$. By induction, there is a set of at most $|U'| = |U| - 1$ vertex disjoint paths in T' each of which ends at a leaf of T' and each vertex in U' lies on exactly one path. Adding the (trivial) path consisting of just y to this set of paths, we obtain a set of vertex disjoint paths in T each ending at a leaf of T and each vertex in U lying on exactly one path. This completes the proof for (ii). \square

Proof of equivalence of (I)–(V) in Theorem 2.1. We establish the following implications between the stated conditions on \mathcal{N} : (I) implies (II), (II) implies (III), (III) implies (IV), and (IV) implies (I), which together show the equivalence of (I)–(IV). We then show that (III) implies (V) and (V) implies (II).

(I) \Rightarrow (II). Suppose that \mathcal{N} is tree-based, and T is a base tree of \mathcal{N} . Let $U = V - X$. Then, by Lemma 3.1, there is a collection of vertex disjoint paths in T each ending at a vertex in X and each vertex in \mathcal{N} lying on exactly one path. Choosing $\mathcal{A} = X$, the vertex sets of these paths form the blocks of the required partition Π of V .

(II) \Rightarrow (III). Suppose that Π is a partition of V with the property that each block in Π is the vertex set of a path in \mathcal{N} ending at a leaf in X . Let U be a subset of vertices of \mathcal{N} . Then Π provides a set of vertex disjoint paths each ending at a leaf in X and with each vertex in U on exactly one path.

(III) \Rightarrow (IV). For this implication, we prove the contrapositive. Suppose that property (IV) is false for \mathcal{N} . Then there exist subsets U_1 and U_2 of V with $|U_1| > |U_2|$ that satisfy the two traversal conditions (i) and (ii). We show $U = U_1$ fails to satisfy property (III). First observe that if P is a path in \mathcal{N} ending at X , then P contains at least as many vertices of U_2 as U_1 . To see this, observe that because of the traversal conditions (i) and (ii), as we move along P , we alternate between vertices in U_1 and vertices U_2 . That is, for $\{i, j\} = \{1, 2\}$, if we traverse a vertex in U_i , then the next vertex we traverse in $U_i \cup U_j$ is a vertex in U_j . Moreover, for each vertex in U_1 on P , there is a subsequent vertex in U_2

on P . Hence there are at least as many vertices of U_2 as U_1 in P . Thus any set of vertex disjoint paths in \mathcal{N} each ending at a leaf in X collectively contains at least as many vertices in U_2 as U_1 . But then it is not possible for such a set of paths to collectively contain all the vertices in U_1 since $|U_2| < |U_1|$. So $U = U_1$ violates property (III).

(IV) \Rightarrow (I). Again, we prove the contrapositive. Suppose that \mathcal{N} is not tree-based. Then, by [Theorem 1.1](#), there is a maximal path in $\mathcal{Z}_{\mathcal{N}}$ that starts and ends in R . Writing this path as $r_1 t_1 r_2 \cdots t_{k-1} r_k$, let q be the parent of r_1 that is not t_1 and let q' be the parent of r_k that is not t_{k-1} . Since the path is maximal, both q and q' are reticulations of \mathcal{N} . Let $U_1 = \{q, t_1, t_2, \dots, t_{k-1}, q'\}$ and $U_2 = \{r_1, r_2, \dots, r_k\}$. These sets have the following properties:

- (1) $|U_1| > |U_2|$,
- (2) U_1 is the set of all parents of all vertices in U_2 , and
- (3) U_2 is the set of all children of all vertices in U_1 .

Now (2) implies that every path from a vertex in U_1 to a vertex in X traverses a vertex in U_2 . Furthermore, (2) and (3) imply that, if there is a path from a vertex in U_1 to another vertex in U_1 , then this path traverses a vertex in U_2 . Similarly, if there is a path from a vertex in U_2 to another vertex in U_2 , then this path traverses a vertex in U_1 . It follows that U_1 and U_2 provide an instance for which property (IV) fails.

(III) \Rightarrow (V). Taking $U = V$ in property (III) immediately gives a path system satisfying property (V).

(V) \Rightarrow (II). Suppose P is a set of paths satisfying property (V). Taking $\mathcal{A} = X$ and $\Pi = P$ gives an antichain and partition of V into $|\mathcal{A}|$ chains that satisfies property (II). \square

Proof of [Theorem 2.2](#). The ‘only if’ direction holds even without the temporal condition, and was established in [\[5\]](#). It also follows from the equivalence of (I) and (III) in [Theorem 2.1](#). For the ‘if’ direction, suppose that \mathcal{N} is temporal but not tree-based. We will show that \mathcal{N} does not possess the antichain-to-leaf property.

Since \mathcal{N} is not tree-based, it follows by [Theorem 1.1](#) that the bipartite graph $\mathcal{Z}_{\mathcal{N}}$ contains a maximal path

$$r_1 t_1 r_2 \cdots t_{k-1} r_k$$

that starts and ends in R . If $k = 1$, then both parents, q and q' say, of r_1 are reticulations, in which case $U = \{q, q'\}$ is an antichain in \mathcal{N} that violates the antichain-to-leaf property. Thus we may assume that $k \geq 2$.

Let q be the parent of r_1 that is not t_1 and let q' be the parent of r_k that is not t_{k-1} . Since the path is maximal, both q and q' are reticulations in \mathcal{N} . Let $U = \{q, t_1, t_2, \dots, t_{k-1}, q'\}$. Since \mathcal{N} is temporal, there is a temporal map λ for \mathcal{N} which necessarily gives

$$\lambda(q) = \lambda(r_1) = \lambda(t_1) = \lambda(r_2) = \dots = \lambda(t_{k-1}) = \lambda(r_k) = \lambda(q'),$$

and so λ is constant on U . If U is an antichain, then U violates the antichain-to-leaf property, since any set of paths that connects the $k + 1$ vertices in U to the leaves in X need to pass through the k vertices in $\{r_1, r_2, \dots, r_k\}$ and so these paths cannot be disjoint.

Therefore, suppose that U is not an antichain. Then there is a directed path P in \mathcal{N} from a vertex $u \in U$ to another vertex $u' \in U$. Moreover, every edge in P must be a reticulation edge of \mathcal{N} . Otherwise, if P contains a tree edge, then $\lambda(u') > \lambda(u)$, contradicting the constancy of λ on U . In particular, the only possible tree vertex of \mathcal{N} in P is the first vertex. Thus P must include q or q' . If q (resp. q') can be reached by a directed path from a vertex in $\{r_1, r_2, \dots, r_k\}$, denote this vertex by r_q (resp. $r_{q'}$). It is easily checked that $r_q \neq r_{q'}$. Now let

$$U' = \begin{cases} U - \{q\}, & \text{if } r_q \text{ exists;} \\ U - \{q'\}, & \text{if } r_{q'} \text{ exists;} \\ U - \{q, q'\}, & \text{if } r_q \text{ and } r_{q'} \text{ exist.} \end{cases}$$

The set U' is an antichain of size k or $k - 1$. Now any path in \mathcal{N} that connects a vertex in U' with a vertex in X must traverse a vertex in $\{r_1, r_2, \dots, r_k\}$. But if r_q exists, then any path traversing r_q must also traverse r_1 . Similarly, if $r_{q'}$ exists, then any path traversing $r_{q'}$ must also traverse r_{k-1} . In all possibilities for U' , it follows that U' does not satisfy the antichain-to-leaf property. \square

4. Measures of deviation

The concept of being tree-based is an ‘all-or-nothing’ property. In this section, we consider computable indices associated with a phylogenetic network \mathcal{N} which are each zero if and only if \mathcal{N} is tree-based. These indices provide, more generally, some measure of how close an arbitrary phylogenetic network is to being tree-based.

Let $\mathcal{N} = (V, E)$ be a phylogenetic network on X . Consider the operation of adjoining a new leaf y to \mathcal{N} by subdividing an edge of \mathcal{N} with a new vertex, u say, and adding the edge (u, y) . Observe that u is a tree vertex in the resulting network. We refer to this operation as *attaching a new leaf* to \mathcal{N} . The three measures we consider are as follows:

- (i) The minimum number $l(\mathcal{N})$ of leaves in $V \setminus X$ that must be present as leaves in a rooted spanning tree of \mathcal{N} .
- (ii) The minimum number $p(\mathcal{N}) = d(\mathcal{N}) - |X|$, where $d(\mathcal{N})$ is the smallest number of vertex disjoint paths that partition the vertices of \mathcal{N} .
- (iii) The minimum number $t(\mathcal{N})$ of leaves that need to be attached to \mathcal{N} so the resulting network is tree-based.

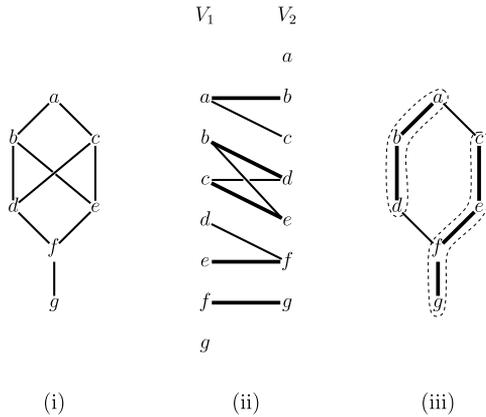


Fig. 3. (i) A phylogenetic network \mathcal{N} on X that is not tree-based, and (ii) the bipartite graph $\mathcal{G}_{\mathcal{N}}$. A maximum-sized matching of $\mathcal{G}_{\mathcal{N}}$ is indicated by the bold edges, and the two corresponding vertex disjoint paths in \mathcal{N} are indicated in (iii). For this example, $u(\mathcal{G}_{\mathcal{N}}) = 2$, and so $p(\mathcal{N}) = u(\mathcal{G}_{\mathcal{N}}) - |X| = 2 - 1 = 1$. A rooted spanning tree with one leaf not in X is obtained from the two paths in (iii) by adding the edge (a, c) .

Each of these measures is non-negative and well defined. To see that (iii) is well defined, attach a new leaf to each reticulation edge in \mathcal{N} . It follows by [Theorem 1.1](#) that the resulting network is tree-based. Moreover, each of these measures equal zero if and only if \mathcal{N} is tree-based. For (ii), this relies on one direction of the equivalence of (I) and (II) in [Theorem 2.1](#).

We will show that each of the measures are computable in time polynomial in the size of \mathcal{N} . Unexpectedly, it turns out that (i), (ii) and (iii) are identical and can be computed by finding a maximum-sized matching in $\mathcal{G}_{\mathcal{N}}$.

Recall that, for a phylogenetic network $N = (V, E)$, the bipartite graph $\mathcal{G}_{\mathcal{N}}$ has vertex bipartition $\{V_1, V_2\}$, where V_1 and V_2 are copies of V , and an edge joining a vertex $u \in V_1$ with a vertex $v \in V_2$ precisely if (u, v) is an edge in \mathcal{N} . Relative to a maximum-sized matching of $\mathcal{G}_{\mathcal{N}}$, let $u(\mathcal{G}_{\mathcal{N}})$ denote the number of unmatched vertices of V_1 . Note that, as each of the elements in X are isolated vertices in V_1 , each of these elements is unmatched regardless of the matching. The first result of this section equates $p(\mathcal{N})$ to the number vertices unmatched by a maximum-sized matching of $\mathcal{G}_{\mathcal{N}}$. An illustration of the proof of this result is given in [Fig. 3](#).

Lemma 4.1. *Let \mathcal{N} be a phylogenetic network on X . Then*

$$p(\mathcal{N}) = u(\mathcal{G}_{\mathcal{N}}) - |X|.$$

Proof. We first show that $p(\mathcal{N}) \leq u(\mathcal{G}_{\mathcal{N}}) - |X|$. Let M be a matching of $\mathcal{G}_{\mathcal{N}}$. Let U_2 denote the set of unmatched vertices in V_2 . For each vertex $u \in U_2$, we recursively construct a directed path P_u in \mathcal{N} as follows. Set $u = u_0$ and initially set $P_u = u_0$. If u_0 is unmatched in V_1 , then terminate the process and set $P_u = u_0$; otherwise, u_0 is matched in V_1 , in which case set $P_u = u_0 u_1$, where $(u_0, u_1) \in M$. If u_1 is unmatched

in V_1 , then terminate the process and set $P_u = u_0 u_1$. Otherwise, u_1 is matched in V_1 , in which case set $P_u = u_0 u_1 u_2$, where $(u_1, u_2) \in M$. Since \mathcal{N} is acyclic, this process eventually terminates with the last vertex, u_k say, added to P_u being unmatched in V_1 .

Repeating this construction for each vertex in U_2 , we eventually obtain a collection $\mathcal{P} = \{P_u : u \in U_2\}$ of directed paths in \mathcal{N} . Since M is a matching, the paths in \mathcal{P} are vertex disjoint. Furthermore, every vertex in \mathcal{N} is on some path in \mathcal{P} . To see this, suppose there is a vertex $v \in V$ not on a path in \mathcal{P} . Clearly, v is matched in V_2 . But then, by reversing the above construction starting at v in V_2 , it is easily seen that v is on such a path. Since each vertex in X is unmatched in V_1 , and noting that the number of paths in \mathcal{P} equates to the number of unmatched vertices in V_2 , and therefore the number of unmatched vertices in V_1 , it follows by choosing M to be of maximum size that

$$p(\mathcal{N}) \leq |\mathcal{P}| - |X| = u(\mathcal{G}_{\mathcal{N}}) - |X|.$$

We next show that $p(\mathcal{N}) \geq u(\mathcal{G}_{\mathcal{N}}) - |X|$. Now let \mathcal{P} be a collection of vertex disjoint paths that partitions the vertices of \mathcal{N} . Let M be the matching of $\mathcal{G}_{\mathcal{N}}$ obtained from \mathcal{P} as follows. The edge $(u, v) \in M$ precisely if u and v are consecutive vertices on some path in \mathcal{P} . Since the paths in \mathcal{P} are vertex disjoint, M is certainly a matching. As every vertex in \mathcal{N} is on some path in \mathcal{P} , the number u_1 of unmatched vertices in V_1 is the number of paths in \mathcal{P} , each such vertex is the last vertex of some path in \mathcal{P} . Thus, by choosing \mathcal{P} to be of minimum size,

$$p(\mathcal{N}) = |\mathcal{P}| - |X| = u_1 - |X| \geq u(\mathcal{G}_{\mathcal{N}}) - |X|.$$

This completes the proof of the lemma. \square

The next corollary establishes the equivalence of (I) and (VI) in [Theorem 2.1](#).

Corollary 4.2. *A phylogenetic network \mathcal{N} on X is tree-based if and only if $\mathcal{G}_{\mathcal{N}}$ has a matching of size $|V| - |X|$.*

Proof. Since the elements of X are always unmatched in V_1 , it follows that $\mathcal{G}_{\mathcal{N}}$ has a matching of size $|V| - |X|$ if and only if $\mathcal{G}_{\mathcal{N}}$ has a maximum-sized matching of this size. In turn, by [Lemma 4.1](#), the latter holds if and only if $p(\mathcal{N}) = u(\mathcal{G}_{\mathcal{N}}) - |X| = 0$. Noting that $p(\mathcal{N}) = 0$ if and only if \mathcal{N} is tree-based completes the proof. \square

We now show that the three measures are identical. An illustration of the first part of the proof of the next theorem is given in [Fig. 3](#).

Theorem 4.3. *Let \mathcal{N} be a phylogenetic network \mathcal{N} on X . Then*

$$l(\mathcal{N}) = p(\mathcal{N}) = t(\mathcal{N}).$$

Proof. We first show that $l(\mathcal{N}) \leq p(\mathcal{N})$. Suppose that Π is the partition of the vertex set V of \mathcal{N} induced by a set of $p(\mathcal{N}) + |X|$ vertex disjoint paths of \mathcal{N} . For each $x \in X$, there is a path ending at x . Consider the paths $\pi_1, \pi_2, \dots, \pi_p$ not ending at an element in X . Note that $p = p(\mathcal{N})$. Since the paths are vertex disjoint and partition V , the set of paths forms a spanning sub-forest of \mathcal{N} . So, by adding, for each path, one edge of \mathcal{N} directed into the starting vertex, we construct a rooted spanning tree T of \mathcal{N} . The leaves of T not in X are precisely the last vertices of the paths $\pi_1, \pi_2, \dots, \pi_p$. Since there are $p(\mathcal{N})$ of these paths, it follows that $l(\mathcal{N}) \leq p(\mathcal{N})$.

We next show that $p(\mathcal{N}) \leq t(\mathcal{N})$. Let \mathcal{N}' be a tree-based network that is obtained from \mathcal{N} by attaching $t(\mathcal{N})$ leaves. Let T be a base tree for \mathcal{N}' , and let U denote the leaf set of T . If we now apply [Lemma 3.1](#) with the same choice of U and T , then T can be partitioned into at most $|U| = t(\mathcal{N}) + |X|$ paths each of which ends at an element in U . Thus $p(\mathcal{N}) \leq t(\mathcal{N})$.

Lastly, we show that $t(\mathcal{N}) \leq l(\mathcal{N})$. Let T be a rooted spanning tree of \mathcal{N} that realises $l(\mathcal{N})$. For each leaf ℓ of T that is not in X , attach a new leaf to an edge directed out of ℓ . If ℓ is a tree vertex of \mathcal{N} , then choose arbitrarily one of the outgoing edges to attach the new leaf. Let \mathcal{N}' denote the resulting phylogenetic network. Since T is a rooted spanning tree of \mathcal{N} , it is easily seen that we can extend T to give a rooted spanning tree of \mathcal{N}' whose leaf set coincides with the leaf set of \mathcal{N}' . Hence \mathcal{N}' is tree-based, and it follows that $t(\mathcal{N}) \leq l(\mathcal{N})$. This completes the proof of the theorem. \square

4.1. Computational complexity and explicit constructions

We now consider the time to compute each of the three measures. Let \mathcal{N} be a phylogenetic network on X and let n denote the total number of vertices in \mathcal{N} . Since each vertex in \mathcal{N} has degree at most three, \mathcal{N} has $O(n)$ edges. By [Theorem 4.3](#),

$$l(\mathcal{N}) = p(\mathcal{N}) = t(\mathcal{N})$$

and, by [Lemma 4.1](#),

$$p(\mathcal{N}) = u(\mathcal{G}_{\mathcal{N}}) - |X|.$$

Thus $l(\mathcal{N})$, $p(\mathcal{N})$, and $t(\mathcal{N})$ can all be computed by finding a maximum-sized matching in $\mathcal{G}_{\mathcal{N}}$. Since $\mathcal{G}_{\mathcal{N}}$ has $2n$ vertices and the same number of edges as \mathcal{N} , we can find such a matching in time $O(n^{3/2})$ [8]. Furthermore, as $\mathcal{G}_{\mathcal{N}}$ can be constructed in time polynomial in n , we can compute $l(\mathcal{N})$, $p(\mathcal{N})$, and $t(\mathcal{N})$ in time polynomial in n .

The proofs of [Lemma 4.1](#) and [Theorem 4.3](#) implicitly establish how one can construct a rooted spanning tree, a system of vertex disjoint paths, and a tree-based network realising $l(\mathcal{N})$, $p(\mathcal{N})$, and $t(\mathcal{N})$. We end this section with explicit algorithms for each of these constructions. Their correctness is omitted as this is essentially done in the proofs of these results. The input to each algorithm is a phylogenetic network \mathcal{N} on X . The

first algorithm constructs a minimum-sized set \mathcal{P} of vertex disjoint paths that partition the vertices of \mathcal{N} .

VERTEX DISJOINT PATHS (\mathcal{N})

1. Construct $\mathcal{G}_{\mathcal{N}}$ and find a maximum-sized matching of $\mathcal{G}_{\mathcal{N}}$.
2. Let U_2 denote the set of unmatched vertices in V_2 .
3. For each $u_0 \in U_2$, find the unique maximal sequence

$$(u_0, u_1), (u_1, u_2), (u_2, u_3), \dots, (u_{k-1}, u_k)$$

of matched edges in $\mathcal{G}_{\mathcal{N}}$ and set P_{u_0} to be the path $u_0 u_1 u_2 \cdots u_k$.

4. Set $\mathcal{P} = \{P_{u_0} : u_0 \in U_2\}$ and return \mathcal{P} .

The second algorithm constructs a rooted spanning tree T of \mathcal{N} that minimises the number of leaves in $V - X$, where V is the vertex set of \mathcal{N} .

ROOTED SPANNING TREE (\mathcal{N})

1. Let \mathcal{P} be the set of vertex disjoint paths returned by a call to VERTEX DISJOINT PATHS (\mathcal{N}).
2. Let π_ρ denote the path in \mathcal{P} traversing the root ρ of \mathcal{N} .
3. For each path $\pi = u_0 u_1 u_2 \cdots u_k$ in $\mathcal{P} - \{\pi_\rho\}$, extend π to

$$\pi' = w u_0 u_1 u_2 \cdots u_k,$$

where (w, u_0) is an edge in \mathcal{N} .

4. Set E_{π_ρ} to be the edge set of π_ρ and, for all $\pi \in \mathcal{P} - \{\pi_\rho\}$, set $E_{\pi'}$ to the edge set of π' .
5. Set $T = (V, E)$, where

$$E = E_{\pi_\rho} \cup \bigcup_{\pi \in \mathcal{P} - \{\pi_\rho\}} E_{\pi'}$$

and return T .

The last algorithm constructs a tree-based network \mathcal{N}' from \mathcal{N} by attaching $t(\mathcal{N})$ leaves.

TREE-BASED NETWORK (\mathcal{N})

1. Let T be the rooted spanning tree returned by a call to ROOTED SPANNING TREE (\mathcal{N}).
2. Let L denote the subset of leaves in T not in X .

3. For each leaf $\ell \in L$, attach a new leaf to an edge directed out of ℓ in \mathcal{N} .
4. Set \mathcal{N}' to be the resulting network and return \mathcal{N}' .

5. Discussion and further questions

In this paper, we have established several new characterisations of tree-based networks based on notions of antichains, path partitions and bipartite matchings. We then applied these results in Section 4 to define and analyse three equivalent and computable measures that quantify how close to being tree-based an arbitrary network is.

There are other ways to quantify the extent to which an arbitrary phylogenetic network deviates from being tree-based; two particular alternative measures are the following:

1. The minimum number of vertices in \mathcal{N} that are required to be absent from any rooted tree that is embedded in \mathcal{N} and has the same leaf set as \mathcal{N} .
2. The minimum number of rooted trees with the same leaf set as \mathcal{N} that are required to be embedded in \mathcal{N} in order that every vertex of \mathcal{N} is present in at least one of the trees.

The first measure is non-negative and equal to zero if and only if \mathcal{N} is tree-based, while the second measure is at least 1, and equal to 1 if and only if \mathcal{N} is tree-based. Determining the computational complexity of computing these two measures seems an interesting topic for future work.

A further project would be to generalise our results concerning the three deviation indices to the larger class of *non-binary* phylogenetic networks, in which vertices are allowed to have in-degree and out-degree greater than two. There are at least two natural ways to extend the definition of tree-based to this more general class of networks, as described and studied in [9].

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