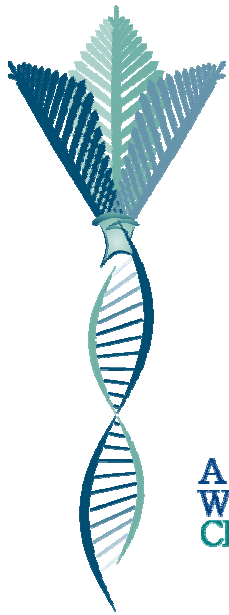

Phylogenetic closure operations and homoplasy-free evolution

IFCS, Chicago, June 2004

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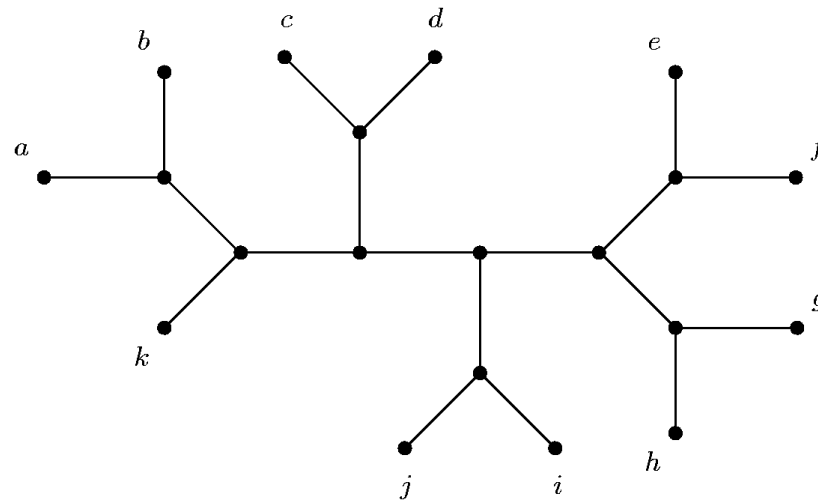


Joint work with...



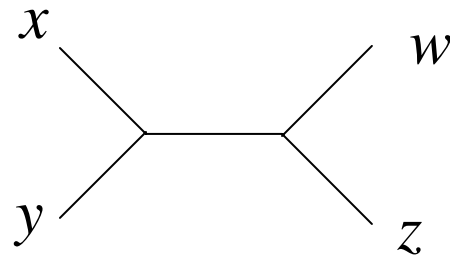
Phylogenetic trees

- [Definition] A **phylogenetic X-tree** is a tree $T=(V,E)$ with a set X of labelled leaves, and all other vertices unlabelled and of degree ≥ 3 .
- If all non-leaf vertices have degree 3 then T is **binary**

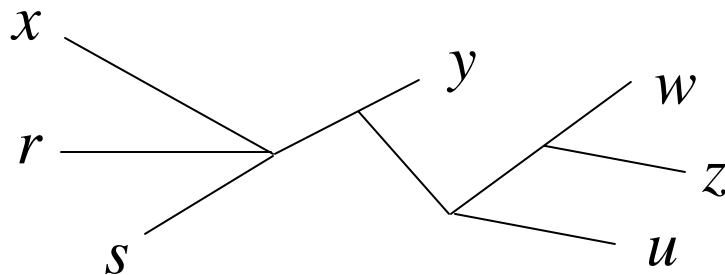


Quartet trees

- A **quartet tree** is a binary phylogenetic tree on 4 leaves (say, x, y, w, z) written $xy|wz$.



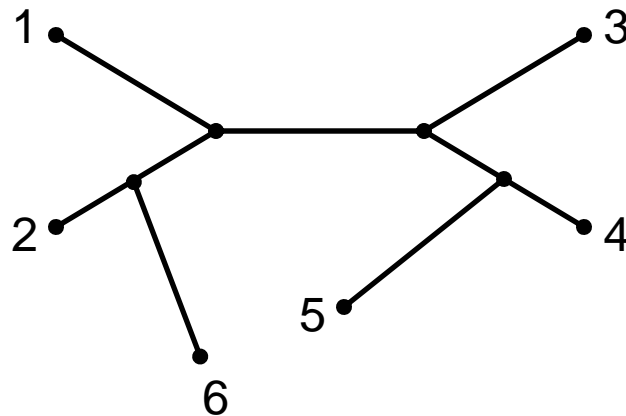
- A phylogenetic X-tree **displays** $xy|wz$ if there is an edge in T whose deletion separates $\{x, y\}$ from $\{w, z\}$



Compatibility

A set Q of quartets is compatible if there is a phylogenetic X -tree T that **displays** each quartet of Q

- **Example:** $Q = \{ 12|34, 13|45, 14|26 \}$



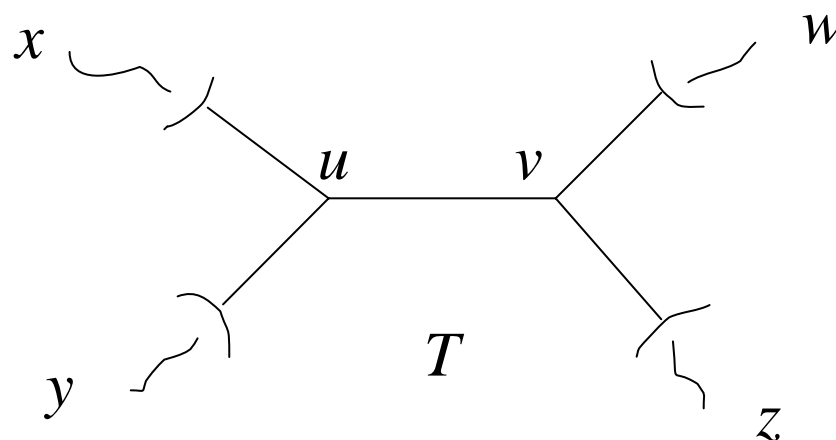
Defining sets

If T is the only phylogenetic X -tree that displays Q (and $X = L(Q)$) then we say Q **defines** T .

- Let $Q(T)$ be the set of **all** quartets displayed by (any) T .
If T is binary, then $Q(T)$ defines T .

A necessary condition for Q to define T

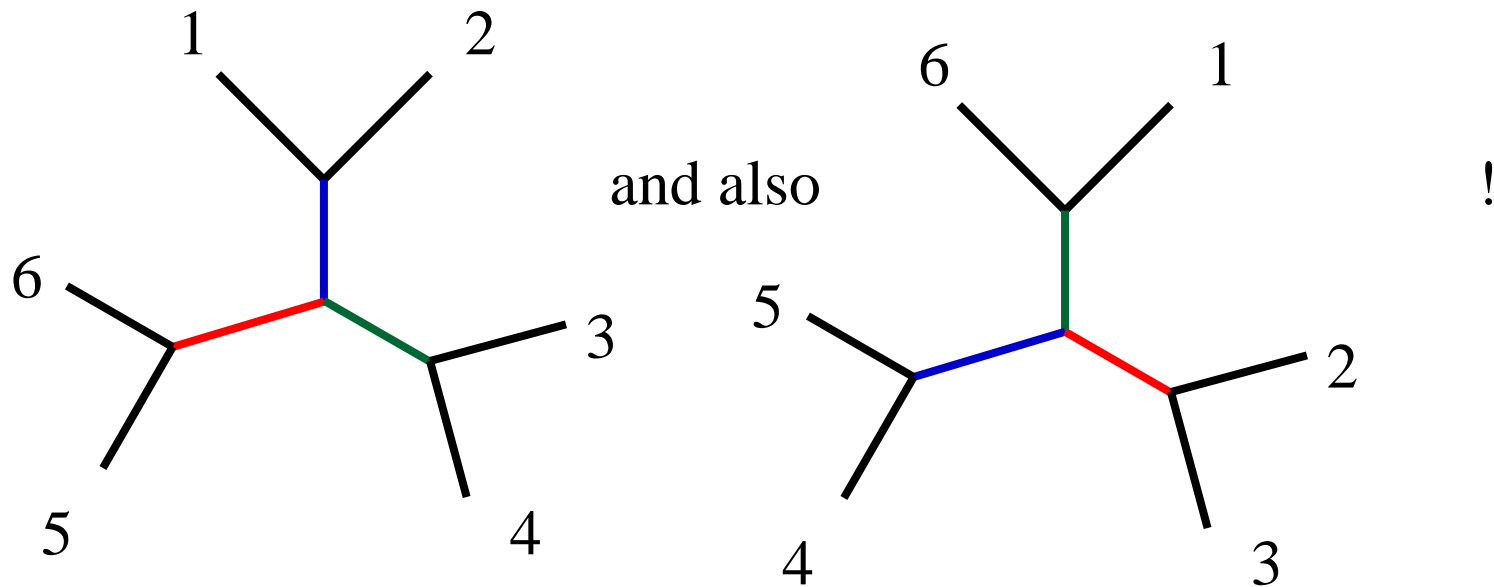
- **Definition:** For a binary phylogenetic tree T , a collection Q of induced quartet trees *distinguishes* an interior edge $\{u,v\}$ of T if there exists a quartet $xy|wz$ in Q that looks like this:



Observation: If Q defines T then T is binary and Q distinguishes every interior edge of T (so $|Q| \geq n-3$).

Warning:

$Q = \{12|45, 56|23, 34|16\}$ distinguishes each interior edge of the tree:



Sufficient condition for Q to define T :

- Suppose Q is compatible and distinguishes every interior edge of a binary phylogenetic X -tree T .

Proposition: If there is an element of X that is a leaf of every tree in Q then Q defines T .

Corollary:

There are subsets of $Q(T)$ that define T of size $|X|-3$.

Character data

Type	States	Transitions
<ul style="list-style-type: none"> Morphology Sequences 	$W(\text{ings}), \neg W, -W$ A, C, G, T	$\neg W \rightarrow W \rightarrow -W$ $x \leftrightarrow y$
<ul style="list-style-type: none"> Gene order Gene content SINEs Oligonucleotides 	$g_1 g_2 g_3 g_4 g_5 g_6 g_7 \dots$ $G = \{g_1, \dots, g_k\}$ $\dots g \dots$ $\dots g_1 g_2 g_3 \dots g_k \dots = 1$	$g_1 g_2 \boxed{g_5 g_4 g_3} g_6 g_7 \dots$ $+g_i / -g_i$ $\dots \rightarrow \dots g \dots \rightarrow \dots ? \dots$ 0 to 1 (once), 1 to 0

Definitions:

- [Character] A character is any function

$$f : X \rightarrow S$$

- [Convexity] Given a character $f : X \rightarrow S$

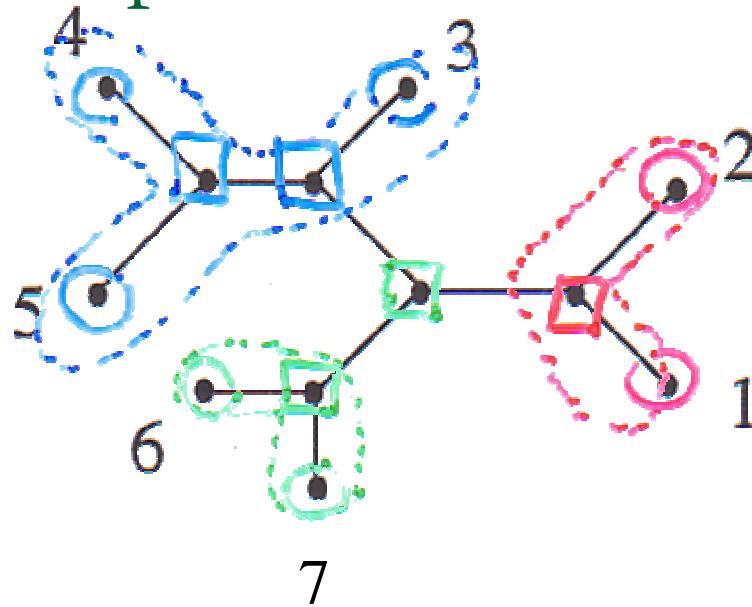
and a phylogenetic X -tree $T=(V,E)$, we say f is **convex on T**

if f extends to $f':V \rightarrow S$

so that $f'|_X = f$

and $\{v \in V : f'(v) = s\}$ is connected for all s in S .

Convexity: example

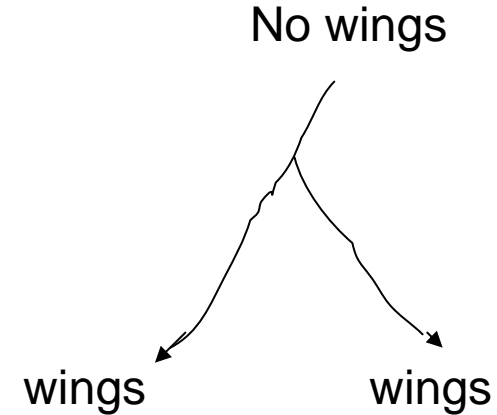
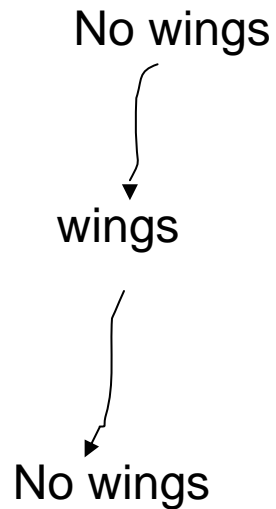


x	1	2	3	4	5	6	7
$f(x)$	•	•	•	•	•	•	•

Biological significance of convexity



- **Lemma:** A character χ is convex on a phylogenetic tree T if and only if χ could have evolved on T (from any root vertex) without any **reversals** or **convergent evolution**.



Relevance to genomics

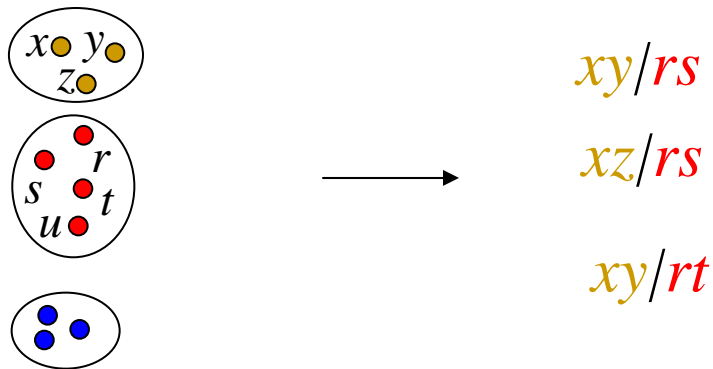
- Eg. gene order rearrangements (n species, L genes, random inversion model)

$$g_1 g_2 \boxed{g_3 g_4 g_5} g_6 g_7 \cdots \quad \Rightarrow \quad g_1 g_2 g_5 g_4 g_3 g_6 g_7 \cdots$$

$$P[h = 0] \geq 1 - \frac{2(2n-3)(n-1)}{L(L-1)}$$

Equivalence of character and quartet compatibility

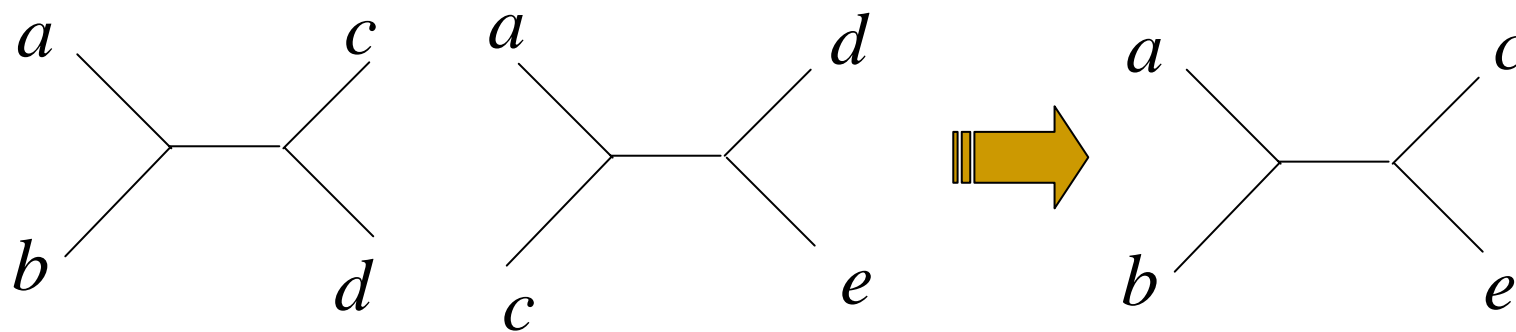
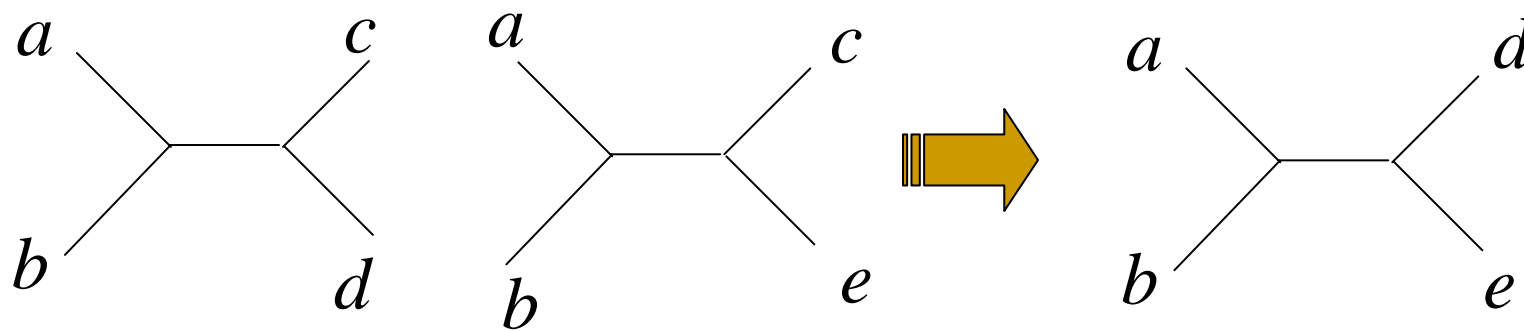
$$C \rightarrow Q(C)$$



Lemma: Each character in C is convex on T if and only if T displays all the quartets in $Q(C)$.

[C is “compatible”, C “defines” T iff $Q(C)$ does]

New quartet trees from old ones



Dyadic rules for quartet trees

(Colonus and Schulze; Dekker)

(Q1): $\{ab|cd, ab|ce\} \vdash ab|de$

(Q2): $\{ab|cd, ac|de\} \vdash ab|ce, ab|de, bc|de.$

Any phylogenetic X -tree that displays the quartet trees on the left of **(Q1)** or **(Q2)** also displays the corresponding quartet tree(s) on the right.

Dyadic quartet closure

$$\mathcal{Q} = \mathcal{Q}_1 \subseteq \mathcal{Q}_2 \subseteq \cdots \subseteq \mathcal{Q}_m = \text{qcl}_\theta(\mathcal{Q})$$

where \mathcal{Q}_{i+1} consists of \mathcal{Q}_i together with all additional quartets that can be obtained from a pair of quartets in \mathcal{Q}_i by applying the rule(s) allowed by θ .

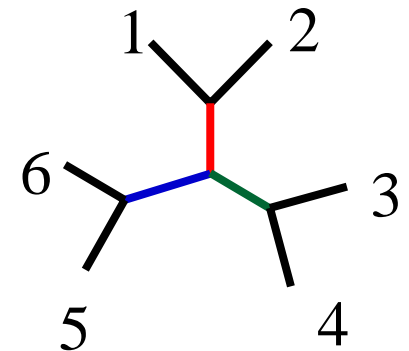
For $\theta \subseteq \{1, 2\}$, let the dyadic quartet closure under rule θ , $\text{qcl}_\theta(\mathcal{Q})$, denote the minimal set of quartet trees that contains \mathcal{Q} and is closed under rule **(Qi)** for each $i \in \theta$.

We denote these closures with: $\text{qcl}_1(\mathcal{Q})$, $\text{qcl}_2(\mathcal{Q})$, $\text{qcl}_{1,2}(\mathcal{Q})$.

Example 1: qcl_2

Definition: If Q distinguishes every interior edge of a binary phylogenetic tree T and we can order Q so that each quartet tree in the ordering introduces precisely one new leaf label, we say Q has a **tight ordering** for T .

Example: $\{12|35, 13|56, 15|34\}$.



Proposition:

If Q has a tight ordering for T , then $\text{qcl}_2(Q) = Q(T)$

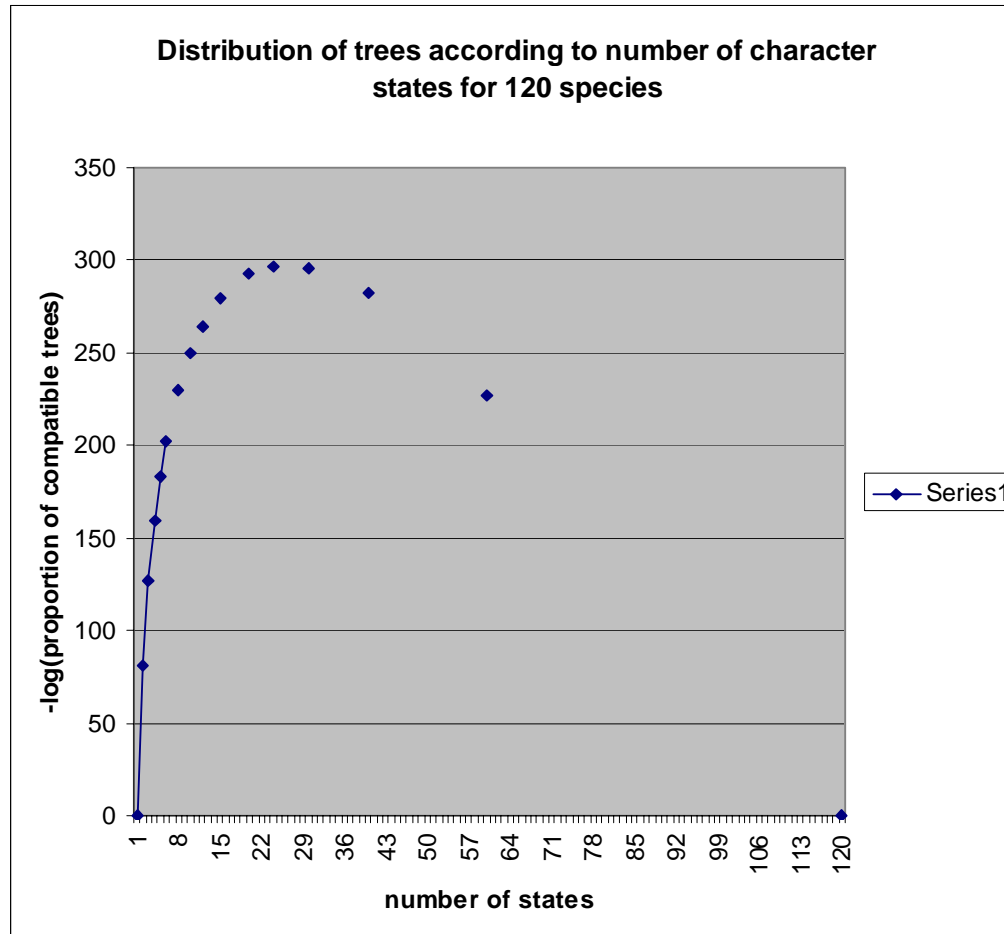
In particular Q defines T .

Application: How many characters are needed to define a binary phylogenetic X -tree?

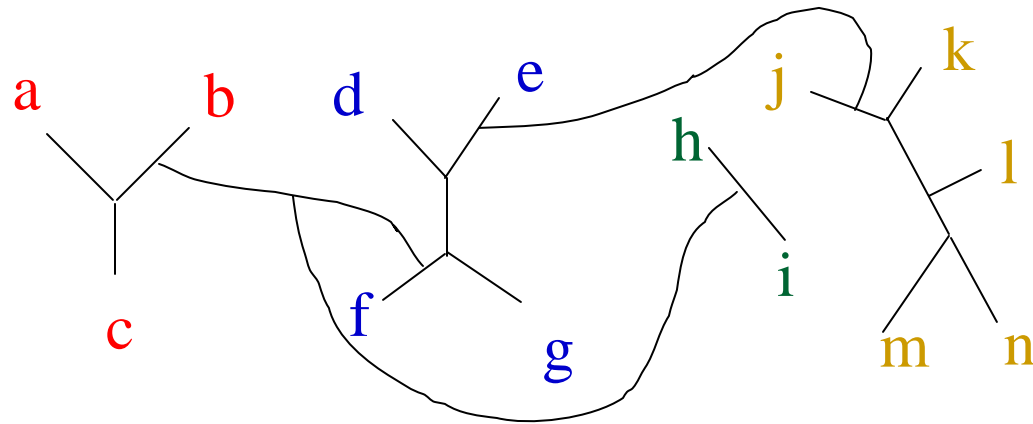
- For binary characters we need $n-3$ ($n=|X|$).
- For r -state characters (r fixed) we need at least
 $(n-3)/(r-1)$
- What if r is not fixed?

(it is not useful to make r too large!)

$$I(\chi) := -\log(\Pr[\chi \text{ is convex on random } T])$$



Where do these numbers come from?

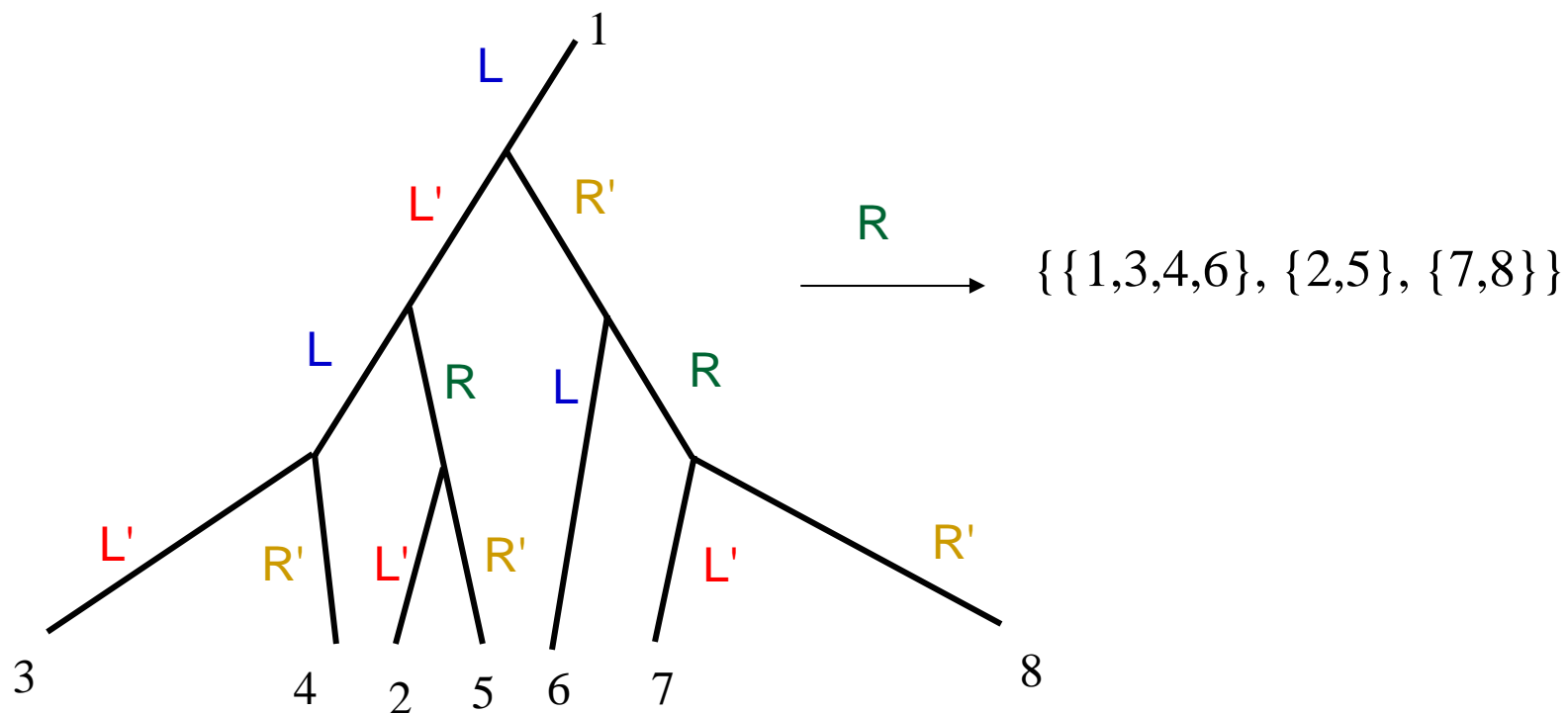


Carter *et al.* (1990); Erdős & Székely (1993).

binary phylogenetic trees with n leaves, $b(n) = 1 \times 3 \times \dots \times (2n - 5)$

of these on which χ is convex =
$$\frac{b(n) \prod_{i=1}^r b(a_i + 1)}{b(n - r + 2)}$$

Edge-colouring a tree by $\mathbb{Z}_2 \times \mathbb{Z}_2$



Theorem (Huber, Moulton, S, 2003)

$Q(C)$ contains a subset with a tight ordering for T .

Thus for any tree there is a set of just *four* characters that defines T .

Distances or characters?

- $d_C(i,j) = \#$ characters in C on which i and j differ

If C is compatible is d_C tree-like?

C binary – yes.

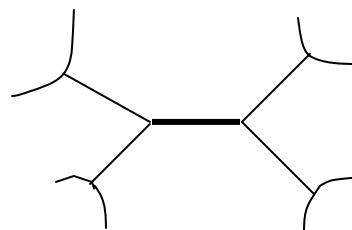
C non-binary no.

Theorem [Huson and S, 2003]:

- For any two trees T_1, T_2 there is a set of multi-state characters C such that
- C defines T_1 (i.e. C homoplasy-free only on T_1) yet
- d_C is tree-like (and ultrametric!), but only on T_2 .

Application 2: “Short” quartets

■ $Q_{\text{short}}(T)$



■ **Theorem** (Erdős *et al.* 1997)

$Q_{\text{short}}(T)$ contains a subset that has a tight ordering for T
(and so $\text{qcl}_2(Q_{\text{short}}(T)) = Q(T)$).

■ The number of characters required to reconstruct (wp $> 1 - \varepsilon$) a binary phylogenetic tree with n leaves from binary characters generated under a finite Markov process is (for almost all trees) at most

$$k \geq \frac{c_\varepsilon (\log(n))^{d(p)}}{a^2}$$

A further application involving qcl_2 :

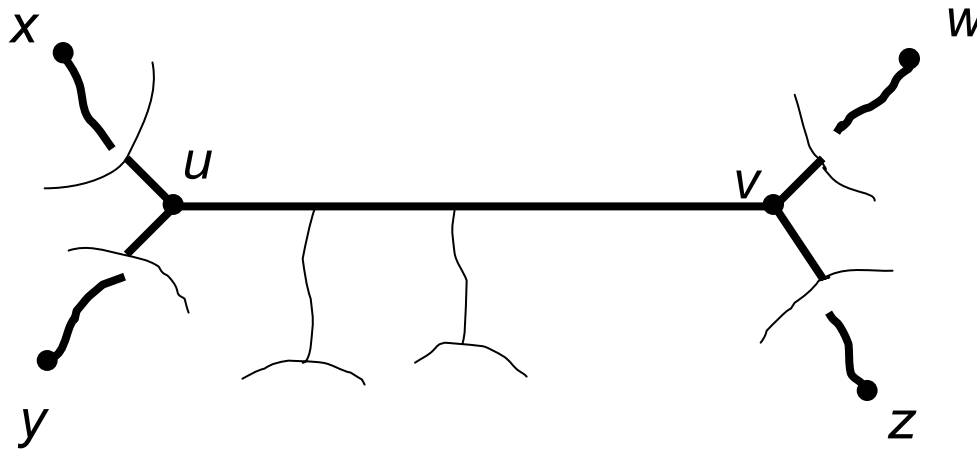
We say Q is **excess-free** if $|L(Q)| - 3|Q| = 0$.

- **Proposition:** Suppose a subset Q of $Q(T)$ contains an excess-free subset Q_0 that defines T . Then $qcl_2(Q) = Q(T)$.
- **Why?** Let us say a set Q of quartet trees is “good” if (i) Q defines a phylogenetic tree, and (ii) $exc(Q) = 0$.

Theorem [Bocker, Dress 1999] Any good set of (≥ 2) quartets is the disjoint union of precisely two good sets.

Example 2: qcl_1

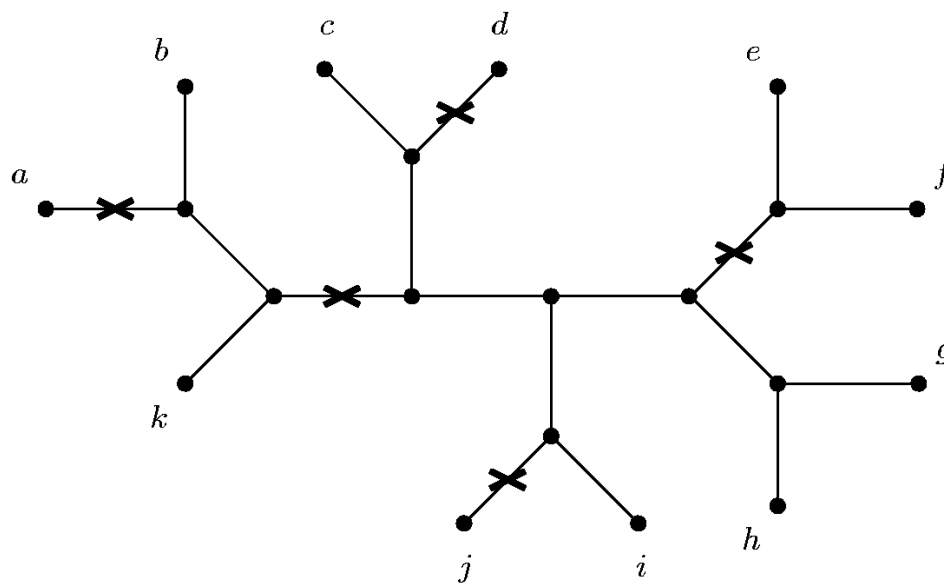
- **Definition:** For a binary phylogenetic tree T , a collection Q of displayed quartet trees is a *generous cover* for T if for all pairs u, v of interior vertices of T , we have a quartet $xy|wz$ in Q that looks like this:



Theorem (Dezulian + S, 2003): If Q is a generous cover for T , then $qcl_1(T) = Q(T)$. Thus Q defines T .

Application: the random cluster model

Random process on a phylogenetic tree \mathcal{T} . Independently cutting edges with probability $p(e)$ generates, by connectivity, random characters on \mathcal{T} .



Cutting the marked edges yields the character $\{a|bk|cghi|d|ef|j\}$.

Reconstructing T from k independent characters (bounds and phase transition)

- **Theorem** (Mossel and S, 2003) For random cluster model, if
- $0 < a \leq p(e) \leq p < 0.5$, every binary phylogenetic tree with n leaves can be reconstructed with probability at least $1 - \varepsilon$ from k indep. characters if

$$k \geq \frac{c_{p,\varepsilon} \log(n)}{a}$$

- A fast (polynomial-time) algorithm to reconstruct T from the characters.
- Proof uses generous cover result. Doesn't require i.i.d.
- Lower bound: $\log(n)$ needed (not trivial) and polynomial (n) if $p > 0.5$.
- *cf.* finite-state $k \geq \frac{c_\varepsilon (\log(n))^{d(p)}}{a^2}$

Relevance to finite-state space?

Corollary:

Random walk on group with generating set of size

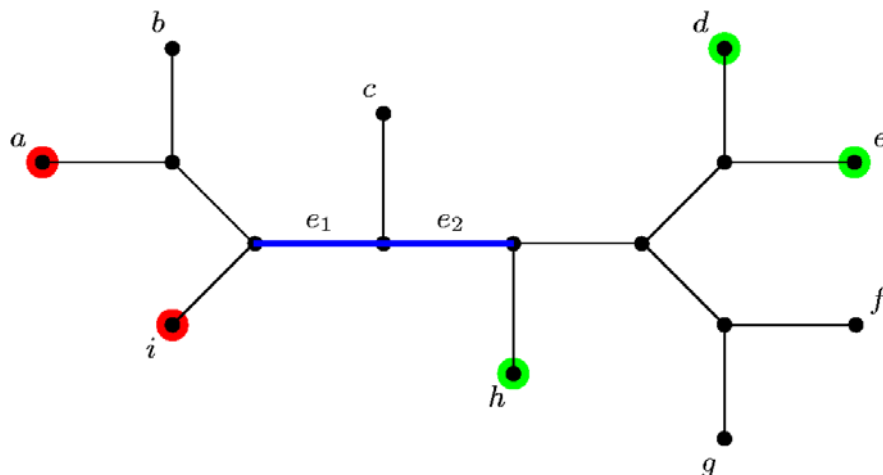
$$d \geq c_\varepsilon n^2 \log(n)$$

then T can be reconstructed w.p. $> 1-\varepsilon$ with a $\Theta(\log(n))$ number of characters

Application 3: qcl_1 , qcl_2 , $qcl_{1,2}$

[Definition] A **partial X-split** $A|B$ is a partition of a subset into two non-empty sets, A, B . $A|B$ is **displayed** by T if we can remove an edge from T to separate A from B .

Example: $\{a, i\} | \{d, e, h\}$ is displayed by T .



Meacham's dyadic rules for splits (1983)

(M1): If $A_1 \cap A_2 \neq \emptyset$ and $B_1 \cap B_2 \neq \emptyset$ then

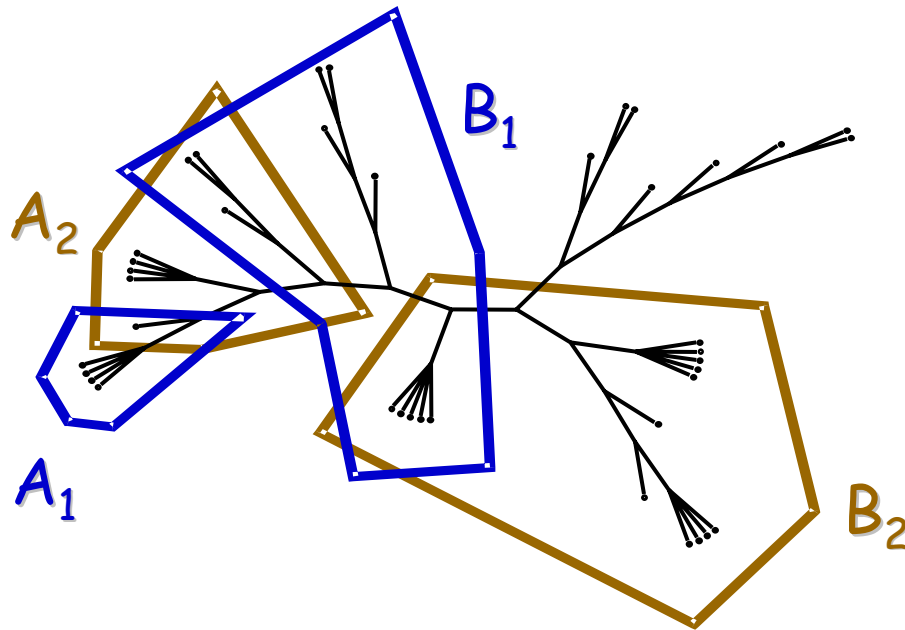
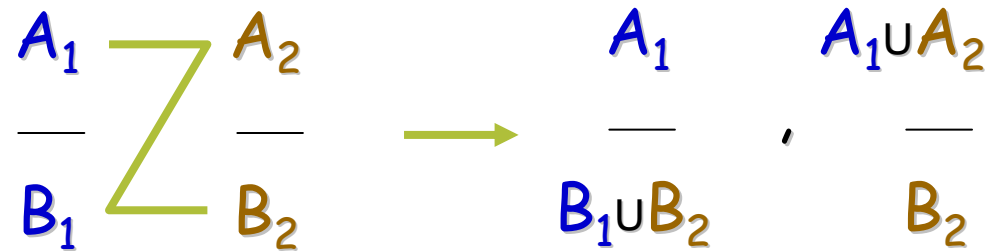
$\{A_1|B_1, A_2|B_2\} \vdash A_1 \cap A_2 | B_1 \cup B_2, A_1 \cup A_2 | B_1 \cap B_2.$

(M2): If $A_1 \cap A_2 \neq \emptyset$ and $B_1 \cap B_2 \neq \emptyset$ and $A_1 \cap B_2 \neq \emptyset$ then

$\{A_1|B_1, A_2|B_2\} \vdash A_2 | B_1 \cup B_2, A_1 \cup A_2 | B_1.$

Any phylogenetic X -tree that displays the partial X -splits on the left of **(M1)** or **(M2)** also displays the corresponding partial X -splits on the right.

Illustration of (M2)



Dyadic split closure

$$\Sigma = \Sigma_1 \subseteq \Sigma_2 \subseteq \cdots \subseteq \Sigma_m = \text{spcl}_\theta(\Sigma)$$

where Σ_{i+1} consists of Σ_i together with all additional splits that can be obtained from a pair of splits in Σ_i by applying the rule(s) allowed by θ .

For $\theta \subseteq \{1, 2\}$, let the dyadic split closure under rule θ , $\text{spcl}_\theta(\Sigma)$, denote the minimal set of splits that contains Σ and is closed under rule **(Mi)** for each $i \in \theta$.

We denote these closures with: $\text{spcl}_1(\Sigma)$, $\text{spcl}_2(\Sigma)$, $\text{spcl}_{1,2}(\Sigma)$.

The (almost) happy marriage

$$\begin{array}{ccc} \Sigma & \xrightarrow{\mathcal{Q}} & \mathcal{Q}(\Sigma) \\ \text{spcl}_\theta \downarrow & & \downarrow \text{qcl}_\theta \\ \text{spcl}_\theta(\Sigma) & \xrightarrow{\mathcal{Q}} & (*) \end{array}$$

?

The (almost) happy marriage

$$\begin{array}{ccc} \Sigma & \xrightarrow{\mathcal{Q}} & \mathcal{Q}(\Sigma) \\ \text{spcl}_\theta \downarrow & & \downarrow \text{qcl}_\theta \\ \text{spcl}_\theta(\Sigma) & \xrightarrow{\mathcal{Q}} & (*) \end{array}$$

Theorem 2.1. *Let Σ be a collection of partial X -splits. Then,*

$$\text{qcl}_\theta(\mathcal{Q}(\Sigma)) = \mathcal{Q}(\text{spcl}_\theta(\Sigma))$$

for $\theta = \{1\}$ and $\theta = \{1, 2\}$. For $\theta = \{2\}$ we have

$$\text{qcl}_\theta(\mathcal{Q}(\Sigma)) \subseteq \mathcal{Q}(\text{spcl}_\theta(\Sigma))$$

-and containment can be strict.

The closure of a set of quartets

For a compatible set \mathcal{Q} of quartet trees, the *closure* $\text{cl}(\mathcal{Q})$ is defined as

$$\text{cl}(\mathcal{Q}) = \bigcap_{T \in \text{co}(\mathcal{Q})} \mathcal{Q}(T)$$

where $\text{co}(\mathcal{Q})$ is the set of phylogenetic trees that display each of the trees in \mathcal{Q} .

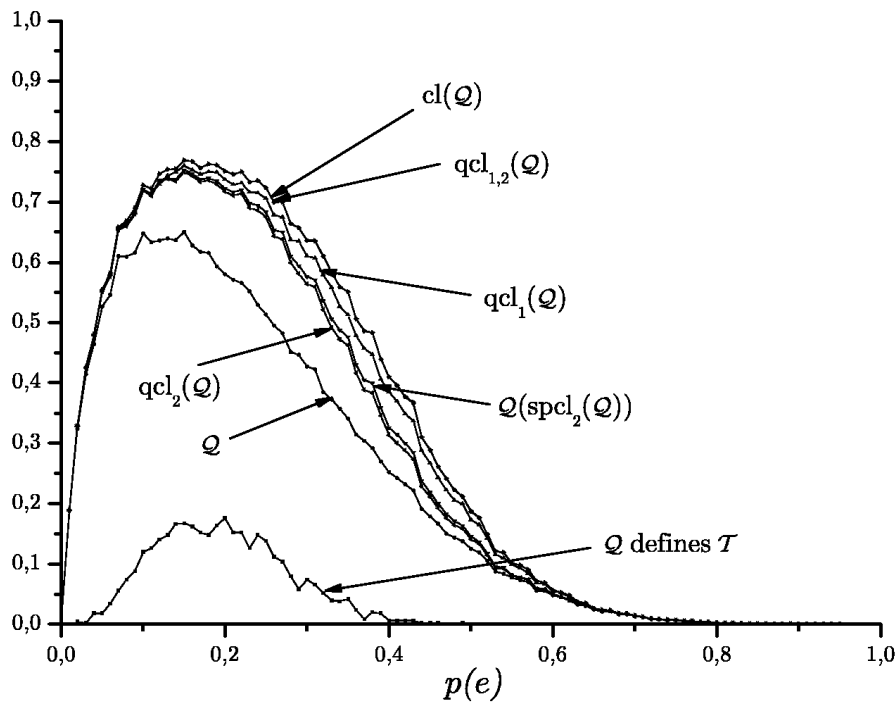
Thus $\text{cl}(\mathcal{Q})$ consists of precisely those quartet trees that are displayed by every phylogenetic tree that displays \mathcal{Q} .

- Rules of order $< p$ (for any fixed p) do not suffice compute $\text{cl}(\mathcal{Q})$.
- There is a set \mathcal{Q} that is incompatible but every strict subset \mathcal{Q}' is compatible and satisfies $\text{cl}(\mathcal{Q}') = \mathcal{Q}'$

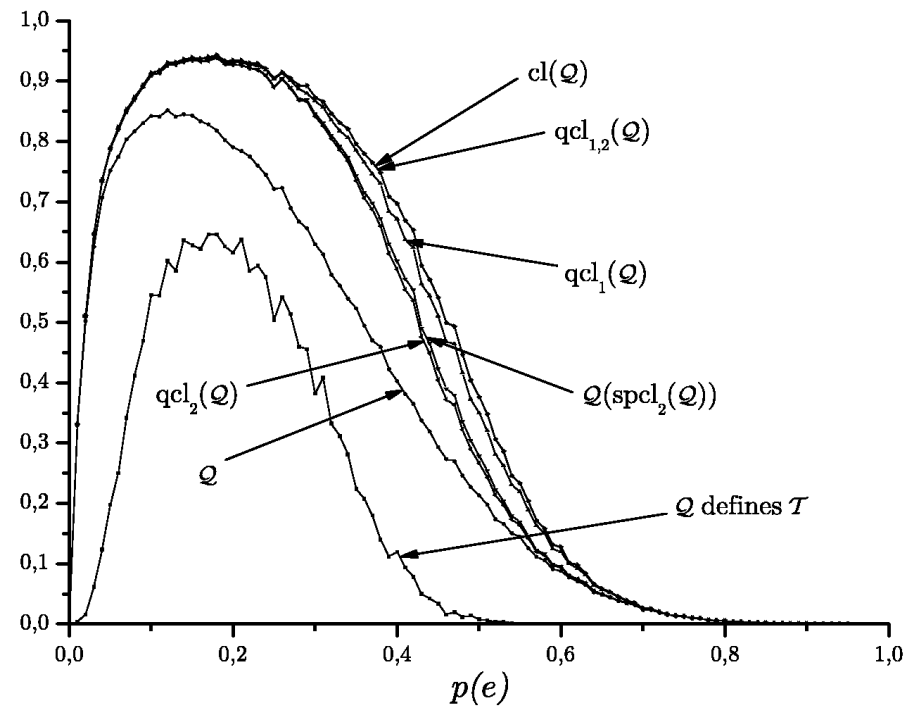
Simulation study

- Random cluster model for binary trees on $n=8$ leaves.
- *Main question:* How much of $\text{cl}(Q)$ does $\text{qcl}_\theta(Q)$ provide? (for $\theta=\{1\}, \{2\}, \{1,2\}$).

Simulations 1 ($n=8, k=16,32$): absolute quartet closure gains



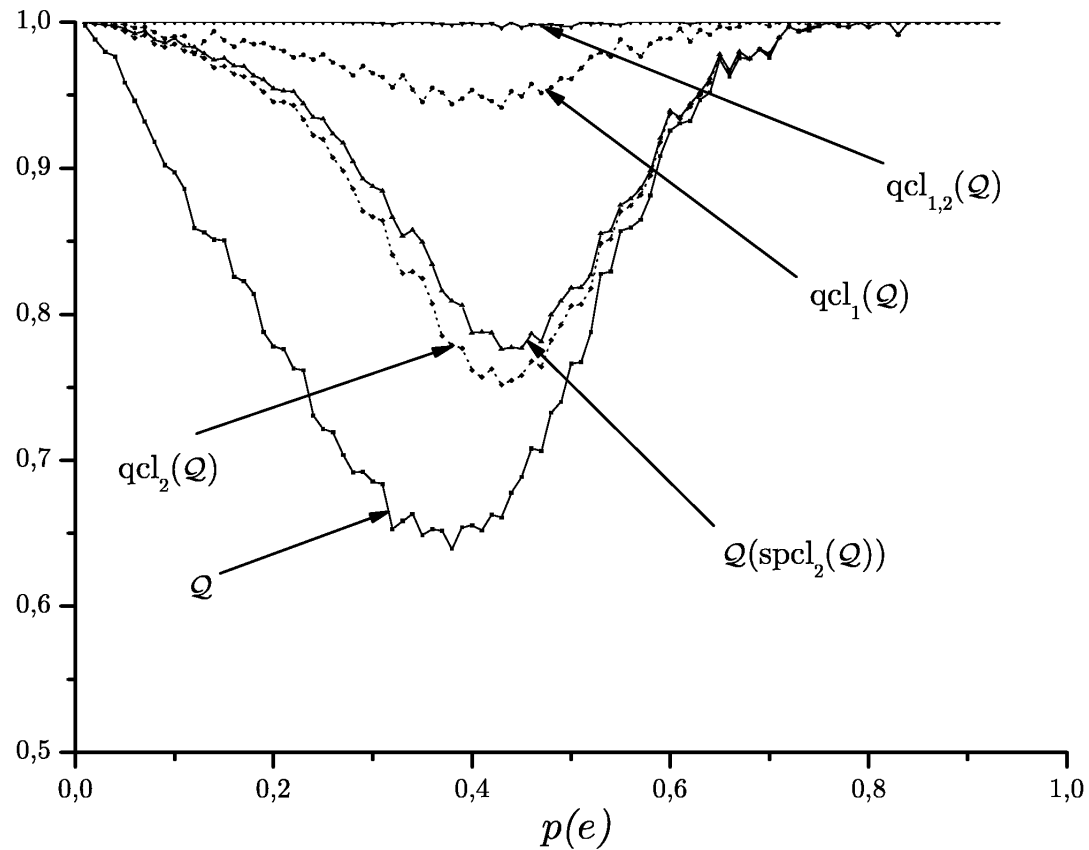
(a) The case $k = 16$.



(b) The case $k = 32$.

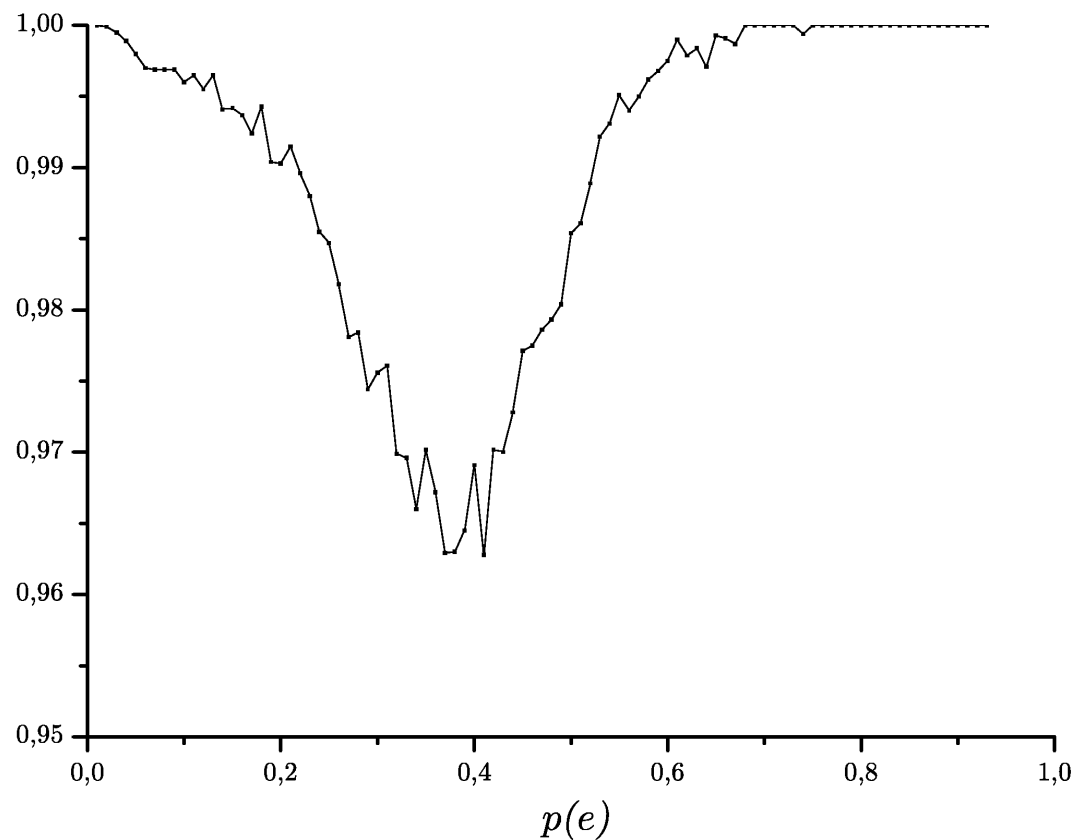
$$\frac{|\mathcal{Q}|}{\binom{n}{4}}, \frac{|\text{qcl}_1(\mathcal{Q})|}{\binom{n}{4}}, \frac{|\text{qcl}_2(\mathcal{Q})|}{\binom{n}{4}}, \frac{|\mathcal{Q}(\text{spcl}_2(\mathcal{Q}))|}{\binom{n}{4}}, \frac{|\text{qcl}_{1,2}(\mathcal{Q})|}{\binom{n}{4}}, \frac{|\text{cl}(\mathcal{Q})|}{\binom{n}{4}}$$

Simulations 2: Relative quartet closure gains



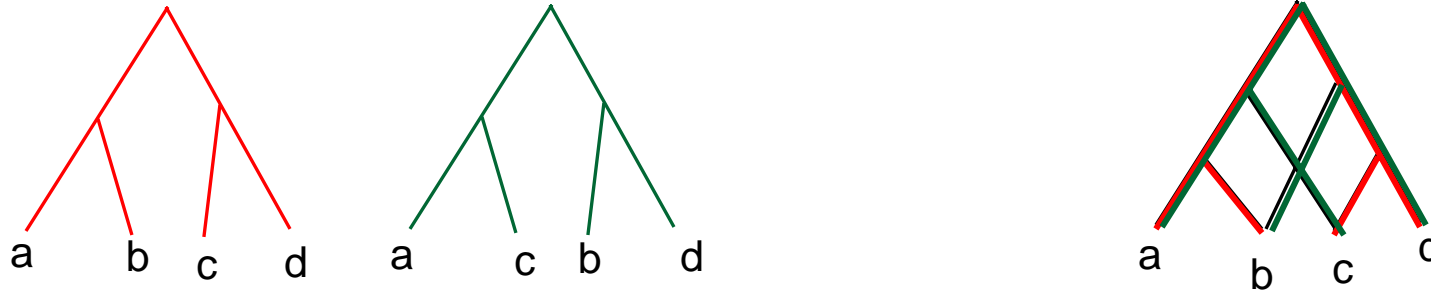
$$\frac{|Q|}{|cl(Q)|}, \frac{|qcl_1(Q)|}{|cl(Q)|}, \frac{|qcl_2(Q)|}{|cl(Q)|}, \frac{|Q(spcl_2(Q))|}{|cl(Q)|}, \frac{|qcl_{1,2}(Q)|}{|cl(Q)|}$$

Simulations 3: dyadic closure of splits and quartets comparison



$\frac{|\text{qcl}_2(\mathcal{Q})|}{|\mathcal{Q}(\text{spcl}_2(\mathcal{Q}))|}$ graphed against $p(e)$

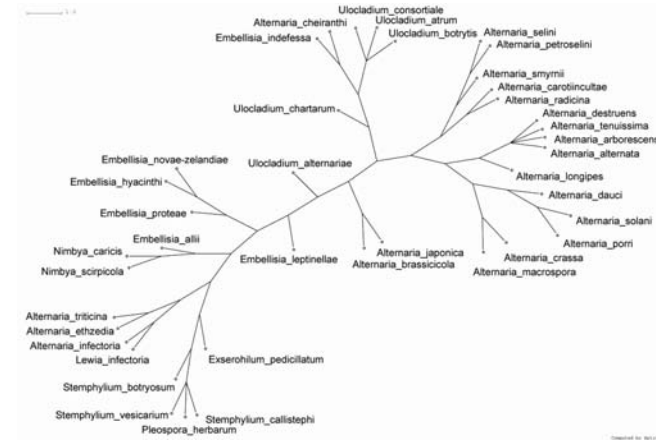
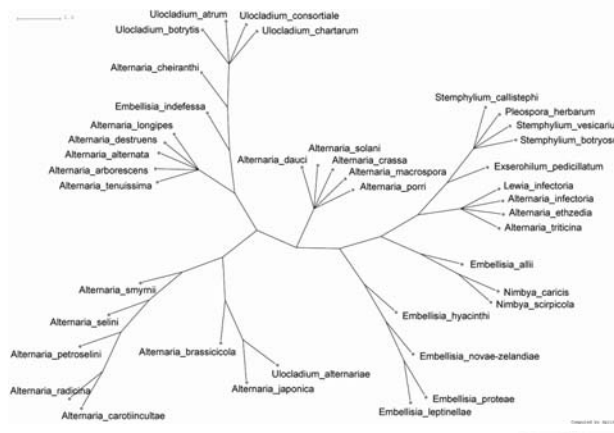
An application of spcl_2

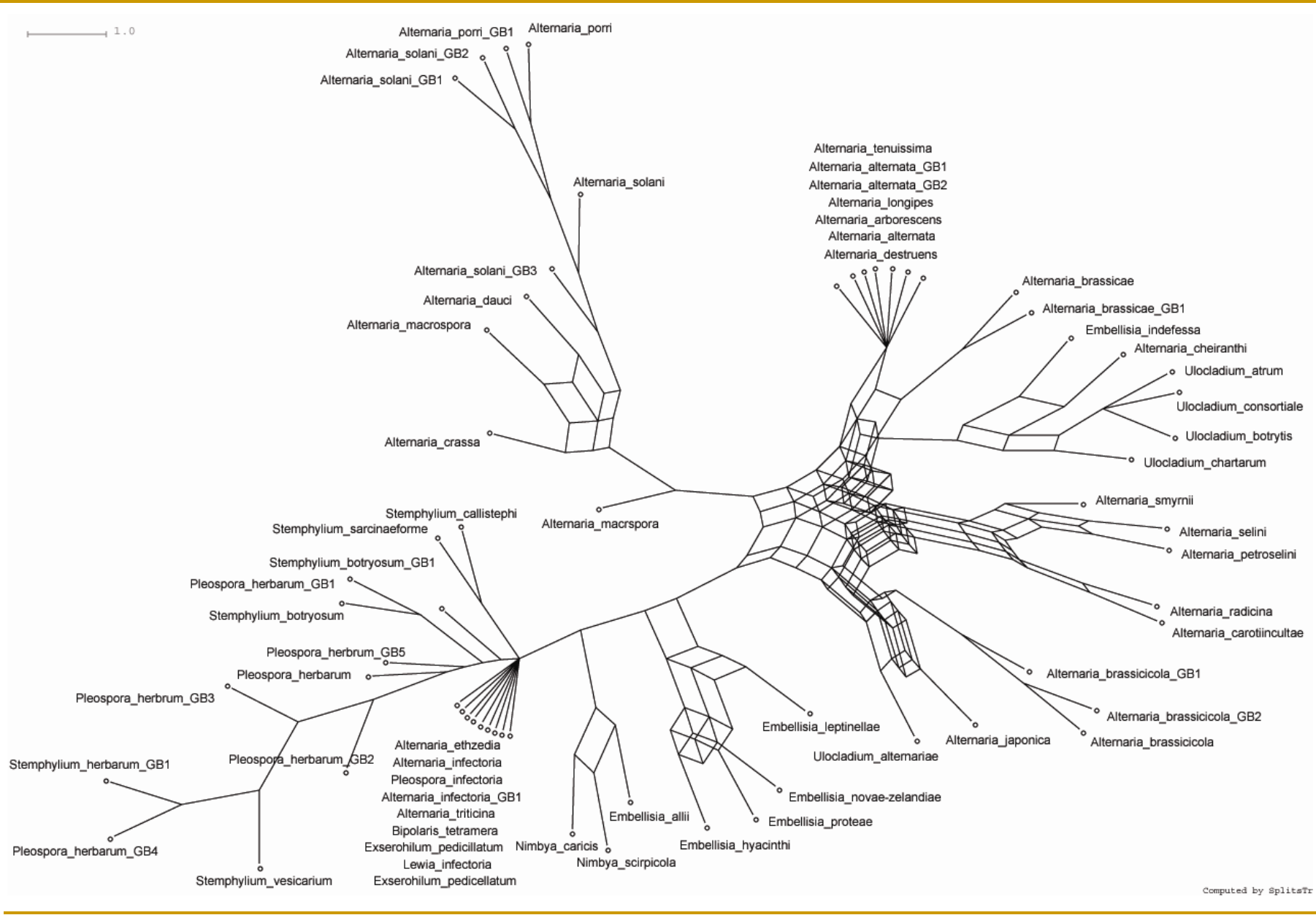


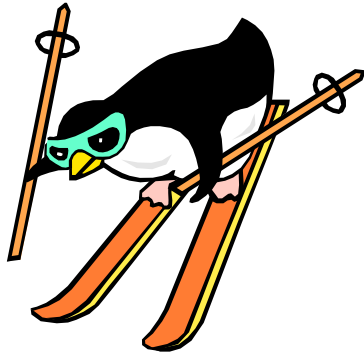
Networks can represent:

- Reticulate evolution (eg. hybrid species)
- Phylogenetic uncertainty (i.e. possible alternative trees)

Approach: Given T_1, \dots, T_k on overlapping sets of species,
let $\Sigma = \Sigma(T_1) \cup \dots \cup \Sigma(T_k)$
construct $\text{spcl}_2(\Sigma)$ and construct the
'splits graph' of the resulting splits that are 'full'.







The end

Further details

- A phase transition for a random cluster model on phylogenetic trees. E. Mossel and M. Steel, *Mathematical Biosciences*, 187 (2004), 189-203.
- Phylogenetic closure operations, and homoplasy-free evolution, T. DeZulian and M. Steel *Proceedings of the International Federation of Classification Societies*, Chicago, 2004.
- Four characters suffice to convexly define a phylogenetic tree. K. Huber, V. Moulton and M. Steel (2003). Submitted.
- How much can evolved characters tell us about the tree that generated them? E. Mossel and M. Steel, Book chapter (Oxford University Press).
- Phylogenetic super-networks from partial trees. D. H. Huson, T. DeZulian, T. Klopper and M. A. Steel, To appear in *WABI 2004*.